

Work with your neighbor. (This will be graded for participation only.)

1. Download <https://www.obagy.com/cs120/LECTURES/sumv1.py> and <https://www.obagy.com/cs120/LECTURES/sumv2.py>.

Add calls to run each version for the values 10,000, 100,000, and 1,000,000. Answer the questions below.

ANS:

What observations do you have about the efficiency of `sumv1.py` vs. `sumv2.py`?

`sumv1.py` is less efficient than `sumv2.py` as the input gets larger

Do you notice a pattern when increasing the input size for `sumv1.py`?

- For `sumv1.py`, as we increase `n`, the running time increases. The running time increases in proportion to `n`.
- For `sumv2.py`, as we increase `n`, the running time stays the same.

2. Consider these two algorithms for searching for a word, `my_word`, in a dictionary (a physical dictionary, which is of course sorted):

- Algo 1: search from the beginning

start at the first word in the dictionary

if the word is not `my_word`, then go to the next word

continue in sequence until `my_word` is found

- Algo 2: start at the middle of the dictionary

if `my_word` is greater than the word in the middle,

start with the middle word and continue from there to the end of the dictionary

if `my_word` is less than the word in the middle,

start with the middle word and search from there to the beginning of the dictionary

ANS:

- a) Which is better, Algo 1 (search from the beginning) or Algo 2 (search from the middle)? Explain your answer.

Most of the time, Algo 2 will be more efficient because for all search words that lie in the second half of the (physical) dictionary, you will eliminate having to check all the words in the first half of the dictionary.

- b) Regardless of which algorithm that you chose, is there ever a scenario where the other one is better? If so, explain the scenario.

For words that lie in the first quarter of the dictionary, Algo 2 would be less efficient because you will first check all the words from the midpoint to the beginning of the second quarter of words before getting to the first quarter of the dictionary.

- c) When considering which is better, what measure are we using?

We are measuring—implicitly counting—the number of comparisons (does `my_word` equal the current word in the dictionary?) that are made before finding the search word.

3. Use the analysis of the `lookup()` example as reference to answer the questions below:

Code	Primitive operations			
<pre>def lookup(str_, list_): for i in range(len(list_)): if str_ == list_[i]: return i return -1</pre>	<table><tr><td>$\left[\begin{array}{ll} \text{len(list_):} & 1 \\ \text{range():} & 1 \\ \text{in :} & 1 \\ \text{for :} & 2 \end{array} \right]$</td><td rowspan="2">each iteration: 9 primitive ops</td></tr><tr><td>$\left[\begin{array}{ll} \text{list_[i] :} & 1 \\ \text{str_ :} & 1 \\ \text{== :} & 1 \\ \text{if :} & 1 \end{array} \right]$</td></tr></table>	$\left[\begin{array}{ll} \text{len(list_):} & 1 \\ \text{range():} & 1 \\ \text{in :} & 1 \\ \text{for :} & 2 \end{array} \right]$	each iteration: 9 primitive ops	$\left[\begin{array}{ll} \text{list_[i] :} & 1 \\ \text{str_ :} & 1 \\ \text{== :} & 1 \\ \text{if :} & 1 \end{array} \right]$
$\left[\begin{array}{ll} \text{len(list_):} & 1 \\ \text{range():} & 1 \\ \text{in :} & 1 \\ \text{for :} & 2 \end{array} \right]$	each iteration: 9 primitive ops			
$\left[\begin{array}{ll} \text{list_[i] :} & 1 \\ \text{str_ :} & 1 \\ \text{== :} & 1 \\ \text{if :} & 1 \end{array} \right]$				
<i>Total primitive ops executed:</i> 1 iteration: 9 ops ∴ n iterations: 9n ops + return at the end: 1 op				

∴ total worst-case running time for a list of length $n = 9n + 1$

- a) What is the worst-case running time of the following code fragment expressed in terms of n ?

```
for i in range(n):  
    k = 2 + 2
```

ANS:

for i in range(n):	n:	1
	range():	1
	in:	1
	for:	2

k = 2 + 2	+:	1
	=:	1

1 iteration: 7 ops
n iterations: 7n ops

Therefore, the worst-case running time for n numbers is 7n

- b) What is the worst-case running time of the following code fragment expressed in terms of n?

```
a = 5
b = 10
for i in range(n):
    x = i * b
    for j in range(n):
        z += b
```

ANS:

a = 5	=:	1	
b = 10	=:	1	total of 2 ops

for i in range(n):	n:	1
	range():	1
	in:	1
	for:	2
x = i * b	i:	1
	b:	1
	*:	1
	=:	1

for j in range(n):	n:	1
	range():	1
	in:	1
	for:	2
z += b	z:	1

	b:	1
	+:	1
	=:	1
at the top		2 ops
1 st loop:	1 iteration:	9 ops
	n iterations:	9n ops
2 nd loop	1 iteration:	9 ops
	n iterations:	9n ops

Therefore, the worst-case running time for n numbers is $18n + 2$